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CS 2102 - DMT1 - FALL 2019 — LUTHER TYCHONIEVICH
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QUIZ 02

PROBLEM 1 *Equivalence rules*

Consider the following equivalence rules:

- the associativity and commutativity of \wedge , \vee , and \oplus
- double negation: $\neg\neg P \equiv P$
- simplification: $P \wedge \perp \equiv \perp$, $P \wedge \top \equiv P$, $P \vee \perp \equiv P$, and $P \vee \top \equiv \top$
- distribution: $A \wedge (B \vee C) \equiv (A \wedge B) \vee (A \wedge C)$ and $A \vee (B \wedge C) \equiv (A \vee B) \wedge (A \vee C)$
- De Morgan: $\neg(A \wedge B) \equiv (\neg A) \vee (\neg B)$ and $\neg(A \vee B) \equiv (\neg A) \wedge (\neg B)$
- definitions: $A \rightarrow B \equiv (\neg A) \vee B$, $(A \leftrightarrow B) \equiv (A \rightarrow B) \wedge (B \rightarrow A)$ and $(A \oplus B) \equiv (A \vee B) \wedge \neg(A \wedge B)$

Prove that $(P \wedge \neg Q) \equiv \neg(P \rightarrow Q)$ by writing out a series of steps, one per line, where the first line is $(P \wedge \neg Q)$, the last line is $\neg(P \rightarrow Q)$, and each line other than the first is an application of **one** of those equivalences to the line above it.

You do not need to name your steps, though doing so might help us grade more easily

$(P \wedge \neg Q)$
 $(\neg\neg P \wedge \neg Q)$
 $\neg(\neg P \vee Q)$
 $\neg(P \rightarrow Q)$

PROBLEM 2 *Prose proof by case analysis*

Write a prose proof of $(P \wedge Q) \rightarrow M \equiv P \rightarrow (Q \rightarrow M)$ by completing the provided template.
Proof. Either P is true or it is false.

Case 1: P is true The expression $(P \wedge Q) \rightarrow M$ in this case

can be simplified to $Q \rightarrow M$ by the equivalence of $\top \wedge Q$ and Q .

The expression $P \rightarrow (Q \rightarrow M)$ in this case

can be simplified to $Q \rightarrow M$ by the equivalence of $\top \rightarrow Q$ and Q .

Because the two are equivalent to the same thing, they are equivalent to each other.

Case 2: P is false The expression $(P \wedge Q) \rightarrow M$ in this case

can be simplified to $\perp \rightarrow M$ by the equivalence of $\perp \wedge Q$ and \perp , which in turn is just \top regardless of the value of M .

The expression $P \rightarrow (Q \rightarrow M)$ in this case

is $\perp \rightarrow (Q \rightarrow M)$, which is \top regardless of the values of Q and M .

Because the two are equivalent to the same thing, they are equivalent to each other.

Since $(P \wedge Q) \rightarrow M \equiv P \rightarrow (Q \rightarrow M)$ is true in both cases, it is true in general. \square