

PROBLEM 1 *Convert to prose*

$S$ : the set of all snakes

$R$ : the set of all rabbits

$E(x, y)$ :  $x$  eats  $y$

$Y(x)$ :  $x$  is yellow

Convert the following to simple, readable English:

1.  $(\exists r \in R, s \in S . E(r, s)) \rightarrow (\neg \forall s \in S . \exists r \in R . E(s, r))$

If any rabbit eats a snake then not all snakes have a rabbit they eat.

2.  $\forall r \in R, s \in S . (Y(s) \rightarrow \neg E(s, r)) \wedge (Y(r) \rightarrow E(r, s))$

Yellow snakes don't eat rabbits, but yellow rabbits eat all the snakes.

3.  $\forall s_1 \in S . \exists s_2 \in S . \forall s_3 \in S . Y(s_1) \rightarrow (\neg E(s_2, s_3) \wedge E(s_1, s_2) \wedge \neg Y(s_2))$

Every yellow snakes eats some non-yellow snake that doesn't eat snakes.

PROBLEM 2 *Primes and factors*

4.            $2^2 \cdot 7$            is the prime factorization of 28

5.            $2^8$            is the prime factorization of 256

6.           31           is the prime factorization of 31

7.            $2^{25} \cdot 7^9$            is the prime factorization of  $4^8 \cdot 14^9$

8.            $\{1, 2, 4, 7, 8\}$            is the set positive 1-digit numbers relatively prime with 15

9.            $\{1, 2, 4, 5, 7, 8\}$            is the set positive 1-digit numbers relatively prime with 81

PROBLEM 3 *Symbolic proof by contradiction*

Write a symbolic proof outline of the the following, using proof-by-contradiction.

10.  $\frac{2}{3} \notin \mathbb{Z}$

Assume  $\frac{2}{3} \in \mathbb{Z}$

$\exists x \in \mathbb{Z} . \frac{2}{3} = x$       definition of set membership

$\frac{2}{3} = x$       existential instantiation

$2 = 3x$       algebra

3 is a factor of 2      fundamental theorem of arithmetic

$\perp$       contradiction

Ergo assumption false      proof by contradiction

$\frac{2}{3} \notin \mathbb{Z}$       conclusion

11.  $\sqrt{2} \notin \mathbb{Q}$

Assume  $\sqrt{2} \in \mathbb{Q}$

$\exists x, y \in \mathbb{Z} . \sqrt{2} = \frac{x}{y}$       definition of rationals

$\sqrt{2} = \frac{x}{y}$       existential instantiation

$2y^2 = x^2$       algebra

2 is a factor of  $2y^2$  with odd multiplicity      fundamental theorem of arithmetic

all factors of  $x^2$  have even multiplicity      fundamental theorem of arithmetic

$\perp$       contradiction

Ergo assumption false      proof by contradiction

$\sqrt{2} \notin \mathbb{Q}$       conclusion

PROBLEM 4 *Prose from symbols*

Write a prose proof that follows the given symbolic proof outlines.

	Assume $\frac{5}{8} \in \mathbb{Z}$	
	$\exists x \in \mathbb{Z} . \frac{5}{8} = x$	definition of set membership
	$\frac{5}{8} = x$	existential instantiation
12.	$5 = 8x$	algebra
	2 is a factor of 5	fundamental theorem of arithmetic
	$\perp$	contradiction
	Ergo assumption false	proof by contradiction
	$\frac{5}{8} \notin \mathbb{Z}$	conclusion

*Proof.*

We proceed by contradiction.

Assume that  $\frac{5}{8}$  is an integer; call that integer  $x$ . This means that  $5 = 8x$ . By the fundamental theorem of algebra, both sides must have the same prime factors; in particular, 2 is a factor of  $8x$ , so it must be a factor of 5; but 2 is not a factor of 5, giving us a contradiction.

Because assume that  $\frac{5}{8}$  is an integer resulted in a contradiction, it must be the case that  $\frac{5}{8}$  is not an integer.  
□

Assume  $\sqrt[3]{4} \in \mathbb{Q}$

$\exists x, y \in \mathbb{Z} . \sqrt[3]{4} = \frac{x}{y} \wedge \gcd(x, y) = 1$  definition of set rationals

$\sqrt[3]{4} = \frac{x}{y}$  existential instantiation

$4y^3 = x^3$  algebra

$\neg(2 \mid x) \vee \neg(2 \mid y)$  because  $\gcd(x, y) = 1$

case 1:  $\neg(2 \mid x)$

$\neg(2 \mid x^3)$

$\perp$

case 2:  $\neg(2 \mid y)$

$(2 \mid x^3)$

$(2 \mid x)$

$(8 \mid x^3)$

$\neg(8 \mid 2y^3)$

$\perp$

$\perp$

Ergo assumption false

$\sqrt[3]{4} \notin \mathbb{Q}$

case analysis

contradiction

proof by contradiction

conclusion

13.

*Proof.*

We proceed by contradiction.

Assume that  $\sqrt[3]{4}$  is a rational number; write that rational in lowest terms as  $\frac{x}{y}$ . This means that  $4y^3 = x^3$ . Because  $\frac{x}{y}$  is in lowest terms, 2 cannot be a factor of both  $x$  and  $y$ ; we thus consider two cases:

**Case: 2 is not a factor of  $x$**  This contradicts the fundamental theorem of arithmetic: because  $4y^3 = x^3$ , 2 must be a factor of  $x^3$  and hence a factor of  $x$  as well.

**Case: 2 is not a factor of  $y$**  By the fundamental theorem of arithmetic, 2 must be a factor of  $x^3$  and hence 8 must be a factor of  $x^3$ ; however, 8 cannot be a factor of  $4y^3$  unless 2 is a factor of  $y$ , resulting in a contradiction.

Because both cases resulted in a contradiction, we have a contradiction in general.

Because assume that  $\sqrt[3]{4}$  is a rational number resulted in a contradiction, it must be the case that  $\sqrt[3]{4}$  is irrational.  $\square$

**PROBLEM 5** *Proof by contradiction*

Prove the following using proof-by-contradiction. You may prove them in prose or in symbols or any readable mix of the two.

14.  $\sqrt{2} \notin \mathbb{Z}$

*Proof.*

We proceed by contradiction.

Assume  $\sqrt{2} \in \mathbb{Z}$ ; let  $x \in \mathbb{Z}$  be the element of  $\mathbb{Z}$  that equals  $\sqrt{2}$ . Thus,  $2 = x^2$ , which means that the prime factorization of  $x$  is  $2^1$ . But a square must have even powers and 1 is not even, which is a contradiction.

Because assuming  $\sqrt{2} \in \mathbb{Z}$  led to a contradiction, it must be the case that  $\sqrt{2} \notin \mathbb{Z}$ .  $\square$

15.  $2^{-1} \notin \mathbb{Z}$

*Proof.*

We proceed by contradiction.

Assume  $2^{-1}$  is an integer; call that integer  $x$ . Then  $2^{-1} = x$ , meaning  $1 = 2x$ . By the fundamental theorem of arithmetic, that means that 2 is a factor of 1, but it is not.

Because assuming  $2^{-1}$  is an integer led to a contradiction, it must be the case that  $2^{-1} \notin \mathbb{Z}$ .  $\square$

16.  $\sqrt{7} \notin \mathbb{Q}$

*Proof.*

We proceed by contradiction.

Assume  $\sqrt{7} \in \mathbb{Q}$ . Then

$$\begin{aligned} \exists x, y \in \mathbb{Z} . \frac{x}{y} &= \sqrt{7} \\ \frac{x}{y} &= \sqrt{7} \\ x &= \sqrt{7}y \\ x^2 &= 7y^2 \end{aligned}$$

But  $x^2$  must have an even number of 7s in its prime factorization and  $7y^2$  must have an odd number, which is a contradiction.

Because assuming  $\sqrt{7} \in \mathbb{Q}$  led to a contradiction, it must be the case that  $\sqrt{7} \notin \mathbb{Q}$ .  $\square$

17.  $3^{1.5} \notin \mathbb{Q}$

*Proof.*

We proceed by contradiction.

Assume  $3^{1.5} \in \mathbb{Q}$ . Let  $\frac{x}{y} = 3^{1.5}$ , where  $x$  and  $y$  are coprime. Then  $\left(\frac{x}{y}\right)^2 = 3^3$  meaning  $x^2 = 3^3y^2$ .  $x^2$  has an even number of 3s in its prime factorization, as does  $y^2$ , meaning  $3^3y^2$  has an odd number. But prime factorization are unique, meaning equal values cannot have differing numbers of 3s in their factorization: thus we have a contradiction.

Because assuming  $3^{1.5} \in \mathbb{Q}$  led to a contradiction, it must be the case that  $3^{1.5} \notin \mathbb{Q}$ .  $\square$

PROBLEM 6 *Additional problems*

18. Prove there are infinitely many prime numbers. Use  $p' = 1 + \prod_{p \in P} p$  where  $P$  is the set of all primes to derive the contradiction (e.g. by showing both that  $p' \in P$  and  $p' \notin P$ ).
19. Prove there are infinitely many integers. Use  $z + 1$  where  $z$  is the largest integer to derive the contradiction.
20. Prove there are infinitely many finite-length strings containing the digits 0 and 1. Use the concatenation of  $s$  and  $s$ , where  $z$  a one of the strings of maximal length, to derive the contradiction.
21. Prove there are infinitely many finite natural numbers. Use  $n + 1$ , where  $n$  is the largest finite natural number, to derive the contradiction.
22. Prove that  $\forall n \in \mathbb{N} . 4 \mid (5^n - 1)$ . Use the well-ordering principle to derive a contradiction by showing that if  $m > 0$  is the smallest  $n$  that makes the expression false, then  $m - 1$  also makes it false. Include a case that shows that the expression holds for  $n = 0$ .
23. Prove that  $\forall n \in \mathbb{Z}^+ . \overline{p_1 \wedge p_2 \wedge \dots \wedge p_n} \equiv \overline{p_1} \vee \overline{p_2} \vee \dots \vee \overline{p_n}$ . Use the well-ordering principle to derive a contradiction by showing that if  $m > 1$  is the smallest  $n$  that makes the expression false, then  $m - 1$  also makes it false. Include a case that shows that the expression holds for  $n = 1$ .
24. Prove there is no smallest positive real number. Use the well-ordering principle to derive a contradiction by showing a smaller positive real number than the smallest positive real. Tools like  $n \div 2$  or  $n \times n$  might help.
25. Prove there is no real number that is closest to, but not the same as,  $x$ . Use the well-ordering principle to derive a contradiction by showing a closer real number than the closest real. Tools like  $\frac{x+y}{2}$  might help.
26. Prove there is no best rational approximation of  $\sqrt{2}$  by showing that, for every approximation  $x$ , the value  $\frac{x}{2} + \frac{1}{x}$  is a better approximation; you may need to a lemma to show that that  $\forall x \in \mathbb{Q} . \frac{x}{2} + \frac{1}{x} \neq x$ .
27. Prove that  $\forall x \in \mathbb{Z} . (x + 1)(x - 1) = x^2 - 1$  without using the distributive law of multiplication. Instead show that it holds for some  $x$  (pick any you wish) and that there's no largest or smallest  $x$  for which it does not hold.
28. Prove that there is no largest two-argument function  $f(x, y)$  that returns  $x + y$  in the programming language of your choice. Do this by showing that if there was a largest program, you can make a larger one that has the same behavior.
29. Prove that there is no most-complicated two-argument function  $f(x, y)$  that returns  $x + y$  in the programming language of your choice, where complication is measured by the number of `if` statements and loops. Do this by showing that if there was a most complicated program, you can make a more complicated one that has the same behavior.
30. Prove that there is no longest-running two-argument function  $f(x, y)$  that returns  $x + y$  in the programming language of your choice. Do this by showing that if there was a most longest-running program, you can make a program that takes longer to execute and has the same behavior.