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CS 2102 - DMT1 - SPRING 2020 — LUTHER TYCHONIEVICH  
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## QUIZ 04

### PROBLEM 1 *Symbolizing*

1. Provide a logic translation for “I’m known by someone so famous that everyone knows them.”  
domain: people

$K(x, y)$ :  $x$  knows  $y$   
 $m$ : Me

$$\exists x . K(x, m) \wedge \forall y . K(y, x)$$

### PROBLEM 2 *Prosimy*

Convert the following proof into a prose proof.

2. Definition: A positive integer is **abundant** if it is smaller than the sum of its factors.

Theorem: There is at least one abundant integer.

Formalism:  $\exists x . A(x)$  (domain: positive integers;  $A(x)$ :  $x$  is abundant)

Proof outline: 12 has factors 1, 2, 3, 4, 6;  $1 + 2 + 3 + 4 + 6 = 16 \vdash A(12)$ ;  $A(12) \vdash \exists x . A(x)$

*Proof.*

Consider the number 12. The factors of 12 are 1, 2, 3, 4, and 6. The sum of those factors is 16, which is greater than 12.

Since the sum of the factors of 12 is greater than 12, 12 is abundant, meaning there is at least one abundant number.

## PROBLEM 3 Complete

Fill in the blanks to complete the following proof by cases that  $P \rightarrow (P \vee Q)$  is a tautology.

*Proof.* 3. Either  $Q$  is true or it is false.

**Case 1:**  $Q$  is true The expression  $P \rightarrow (P \vee Q)$  in this case

4. can be simplified to  $P \rightarrow \top$ ,

which is equivalent to  $\top$ .

**Case 2:**  $Q$  is false The expression  $P \rightarrow (P \vee Q)$  in this case

5. can be simplified to  $P \rightarrow P$ ,

which is equivalent to  $\top$ .

Since  $P \rightarrow (P \vee Q)$  is true in both cases, it is true in general, meaning it is a tautology.  $\square$

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### Symbols

Concept	Java/C	Python	This class	Bitwise	Other
true	true	True	$\top$ or 1	-1	$\top$ , tautology
false	false	False	$\perp$ or 0	0	$\perp$ , contradiction
not $P$	!p	not p	$\neg P$ or $\bar{P}$	$\sim p$	
$P$ and $Q$	p && q	p and q	$P \wedge Q$	p & q	$PQ, P \cdot Q$
$P$ or $Q$	p    q	p or q	$P \vee Q$	p   q	$P + Q$
$P$ xor $Q$	p != q	p != q	$P \oplus Q$	p ^ q	$P \vee Q$
$P$ implies $Q$			$P \rightarrow Q$		$P \supset Q, P \Rightarrow Q$
$P$ iff $Q$	p == q	p == q	$P \leftrightarrow Q$		$P \Leftrightarrow Q, P \text{ xnor } Q$

### Axioms: Equivalence rules

- associativity and commutativity of  $\wedge$ ,  $\vee$ , and  $\oplus$ ; commutativity of  $\leftrightarrow$
- double negation:  $\neg\neg P \equiv P$
- simplification:  $P \wedge \perp \equiv P \wedge \neg P \equiv \perp$ ,  $P \vee \top \equiv P \vee \neg P \equiv \top$ , and  $P \wedge \top \equiv P \vee \perp \equiv P \wedge P \equiv P \vee P \equiv P$
- distribution:  $A \wedge (B \vee C) \equiv (A \wedge B) \vee (A \wedge C)$  and  $A \vee (B \wedge C) \equiv (A \vee B) \wedge (A \vee C)$
- De Morgan:  $\neg(A \wedge B) \equiv (\neg A) \vee (\neg B)$  and  $\neg(A \vee B) \equiv (\neg A) \wedge (\neg B)$
- definitions:  $A \rightarrow B \equiv (\neg A) \vee B$ ,  $(A \leftrightarrow B) \equiv (A \rightarrow B) \wedge (B \rightarrow A)$  and  $(A \oplus B) \equiv (A \vee B) \wedge \neg(A \wedge B)$