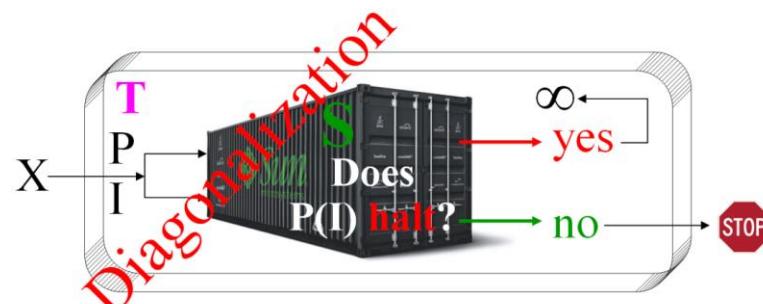


Algorithms

1. Existence

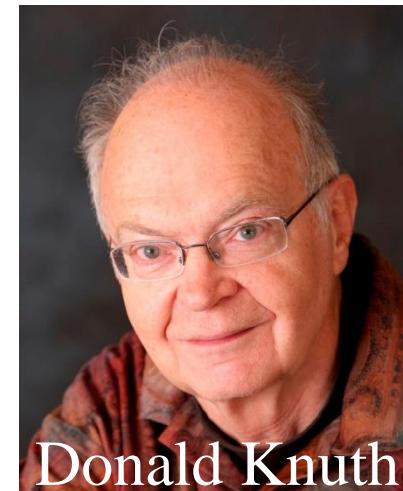
2. Efficiency

- Time
- Space



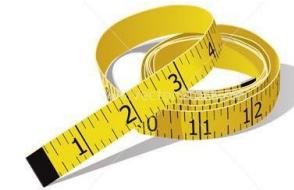
Worst case behavior analysis
as a function of input size

Asymptotic growth: O Ω Θ \circ ω



Donald Knuth

Upper Bounds



Definition: $f(n) = O(g(n))$

$$\Leftrightarrow \exists c, k > 0 \ \exists 0 \leq f(n) \leq c \cdot g(n) \ \forall n > k$$

$$\Leftrightarrow \lim_{n \rightarrow \infty} f(n) / g(n) \text{ exists}$$

$$O(g(n)) = \{f \mid \exists c, k > 0 \ \exists 0 \leq f(n) \leq c \cdot g(n) \ \forall n > k\}$$

“ $f(n)$ is big-O of $g(n)$ ”

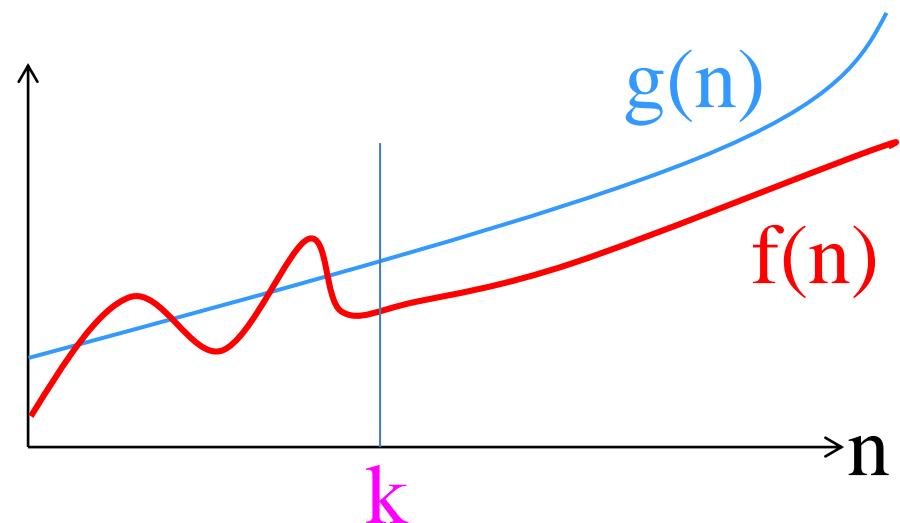
Ex: $n = O(n^2)$

$$33n+17 = O(n)$$

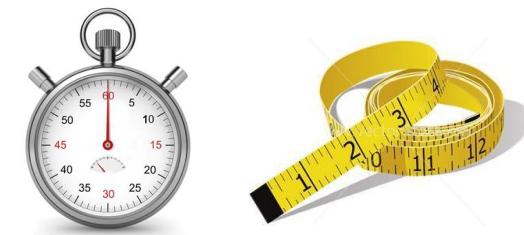
$$n^8+n^7 = O(n^{12})$$

$$n^{100} = O(2^n)$$

$$213 = O(1)$$



Lower Bounds



Definition: $f(n) = \Omega(g(n))$

$$\Leftrightarrow g(n) = O(f(n))$$

$$\Leftrightarrow \lim_{n \rightarrow \infty} g(n) / f(n) \text{ exists}$$

$$\Omega(g(n)) = \{f \mid \exists c, k > 0 \ \exists 0 \leq g(n) \leq c \cdot f(n) \ \forall n > k\}$$

“ $f(n)$ is Omega of $g(n)$ ”

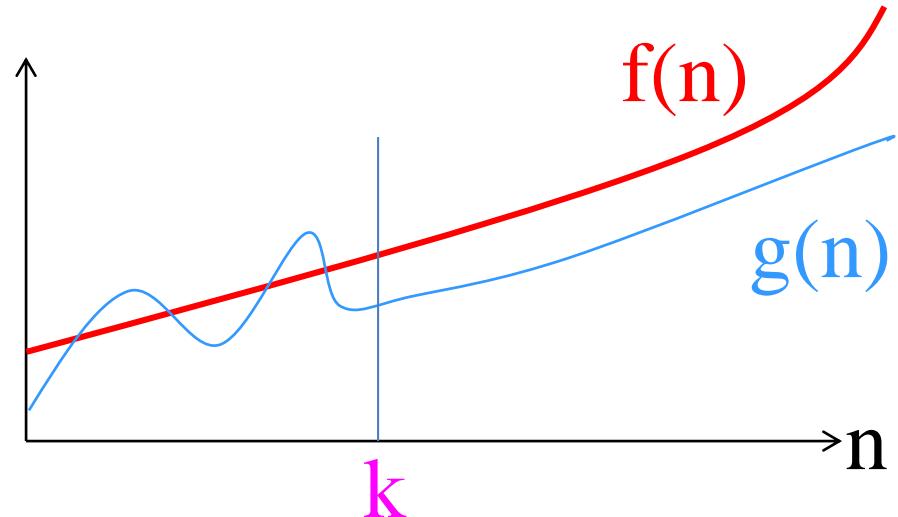
Ex: $100n = \Omega(n)$

$$33n + 17 = \Omega(\log n)$$

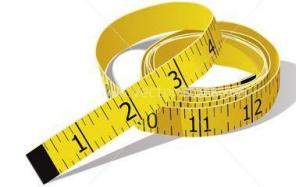
$$n^8 - n^7 = \Omega(n^8)$$

$$213 = \Omega(1/n)$$

$$10^{100} = \Omega(1)$$



Tight Bounds



Definition: $f(n) = \Theta(g(n))$

$\Leftrightarrow f(n) = O(g(n))$ and $g(n) = O(f(n))$

$\Leftrightarrow f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$

$\Leftrightarrow \lim_{n \rightarrow \infty} g(n)/f(n)$ and $\lim_{n \rightarrow \infty} f(n)/g(n)$ exist

“ $f(n)$ is Theta of $g(n)$ ”

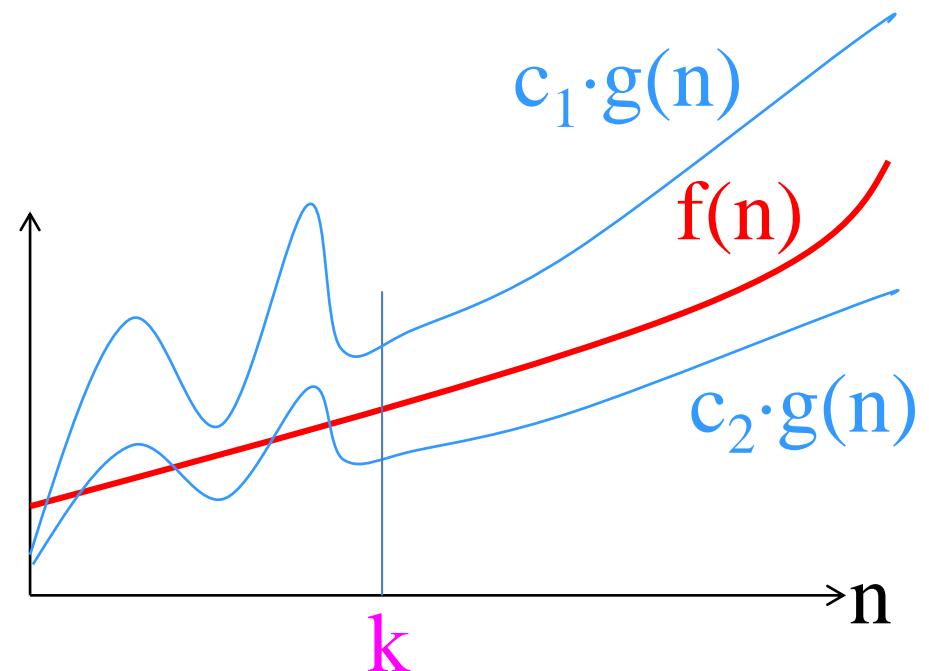
Ex: $99n = \Theta(n)$

$n + \log n = \Theta(n)$

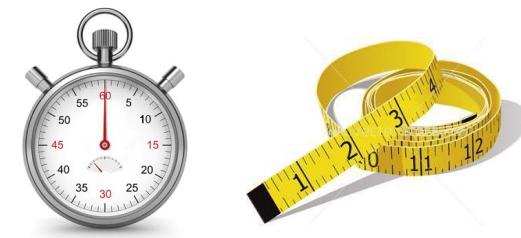
$n^8 - n^7 = \Theta(n^8)$

$n^2 + \cos(n) = \Theta(n^2)$

$213 = \Theta(1)$



Loose Bounds



Definition: $f(n) = o(g_1(n))$

$\Leftrightarrow f(n) = O(g_1(n))$ and $f(n) \neq \Omega(g_1(n))$

“ $f(n)$ is little-o of $g_1(n)$ ”

Definition: $f(n) = \omega(g_2(n))$

$\Leftrightarrow f(n) = \Omega(g_2(n))$ and $f(n) \neq O(g_2(n))$

“ $f(n)$ is little-omega of $g_2(n)$ ”

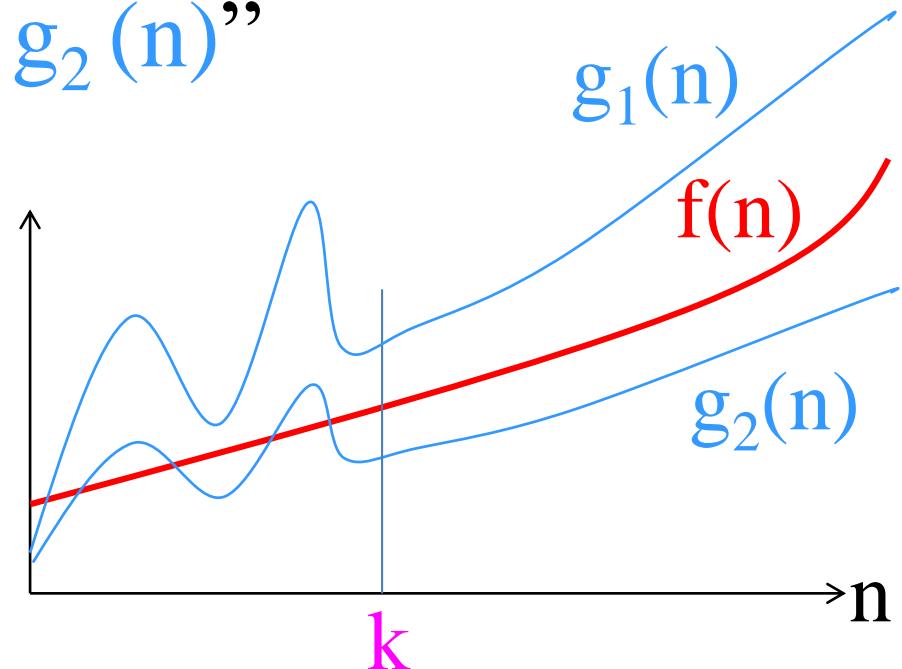
Ex: $8n = o(n \log \log n)$

$n \log n = \omega(n)$

$n^6 = o(n^{6.01})$

$n^2 + \sin(n) = \omega(n)$

$213 = o(\log n)$



Growth Laws



Let $f_1(n) = O(g_1(n))$ and $f_2(n) = O(g_2(n))$

Thm: $f_1(n) + f_2(n) = O(\max(g_1(n), g_2(n)))$

- Sequential code

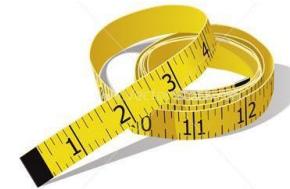
Thm: $f_1(n) \cdot f_2(n) = O(g_1(n) \cdot g_2(n))$

- Nested loops & subroutine calls

Thm: $n^k = O(c^n) \quad \forall c, k > 0$

Ex: $n^{1000} = O(1.001^n)$

Solving Recurrences



$$T(n) = \textcolor{blue}{a} \cdot T(n/\textcolor{green}{b}) + \textcolor{magenta}{f}(n)$$

$a \geq 1$, $b > 1$, and let $\textcolor{red}{c} = \log_{\textcolor{green}{b}} \textcolor{blue}{a}$

Thm: $\textcolor{magenta}{f}(n) = O(n^{\textcolor{red}{c}-\varepsilon})$ for some $\varepsilon > 0 \Rightarrow T(n) = \Theta(n^{\textcolor{red}{c}})$

$\textcolor{magenta}{f}(n) = \Theta(n^{\textcolor{red}{c}}) \Rightarrow T(n) = \Theta(n^{\textcolor{red}{c}} \log n)$

$\textcolor{magenta}{f}(n) = \Omega(n^{\textcolor{red}{c}+\varepsilon})$ some $\varepsilon > 0$ and $\textcolor{blue}{a} \cdot \textcolor{magenta}{f}(n/\textcolor{green}{b}) \leq d \cdot \textcolor{magenta}{f}(n)$
for some $d < 1 \quad \forall n > n_0 \Rightarrow T(n) = \Theta(\textcolor{magenta}{f}(n))$

Ex: $T(n) = 2T(n/2) + n \Rightarrow T(n) = \Theta(n \log n)$

$T(n) = 9T(n/3) + n \Rightarrow T(n) = \Theta(n^2)$

$T(n) = T(2n/3) + 1 \Rightarrow T(n) = \Theta(\log n)$

Stirling's Formula

Factorial: $n! = 1 \cdot 2 \cdot 3 \cdot \dots \cdot (n - 2) \cdot (n - 1) \cdot n$

Theorem: $n! = \sqrt{2\pi n} \cdot \left(\frac{n}{e}\right)^n \cdot \left(1 + \Theta\left(\frac{1}{n}\right)\right)$

where e is Euler's constant = 2.71828...

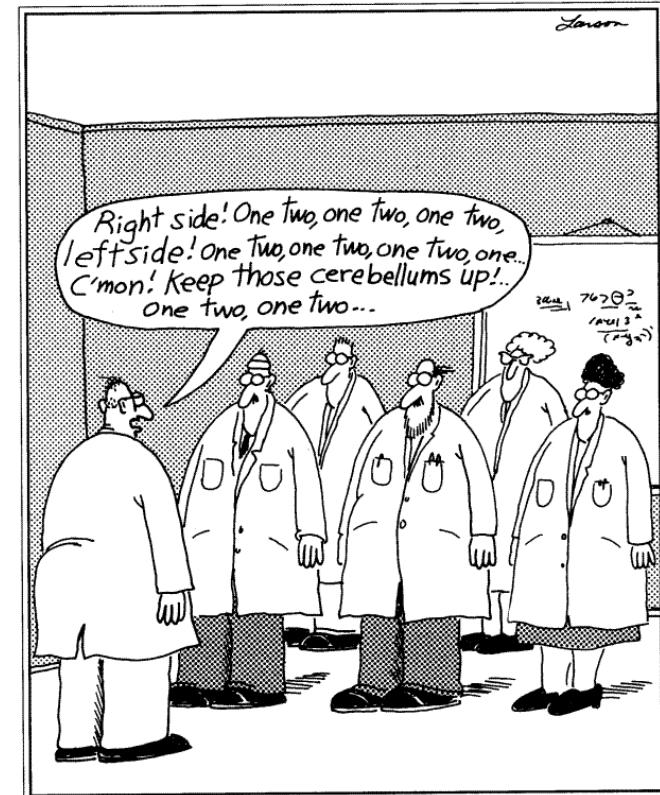
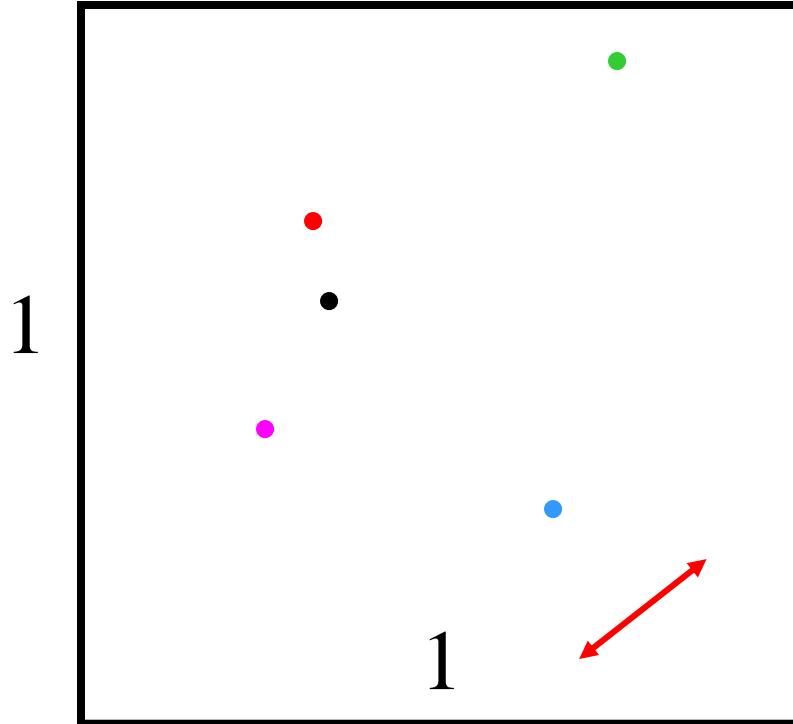
Theorem: $n! \approx \sqrt{2\pi n} \cdot \left(\frac{n}{e}\right)^n$

Corollary: $\log(n!) \approx \log\left(\sqrt{2\pi n} \cdot \left(\frac{n}{e}\right)^n\right) = n \cdot \log\left(\frac{n}{e}\right) + \frac{\log(2\pi n)}{2} = O(n \log n)$

$$\log(n!) = O(n \log n)$$

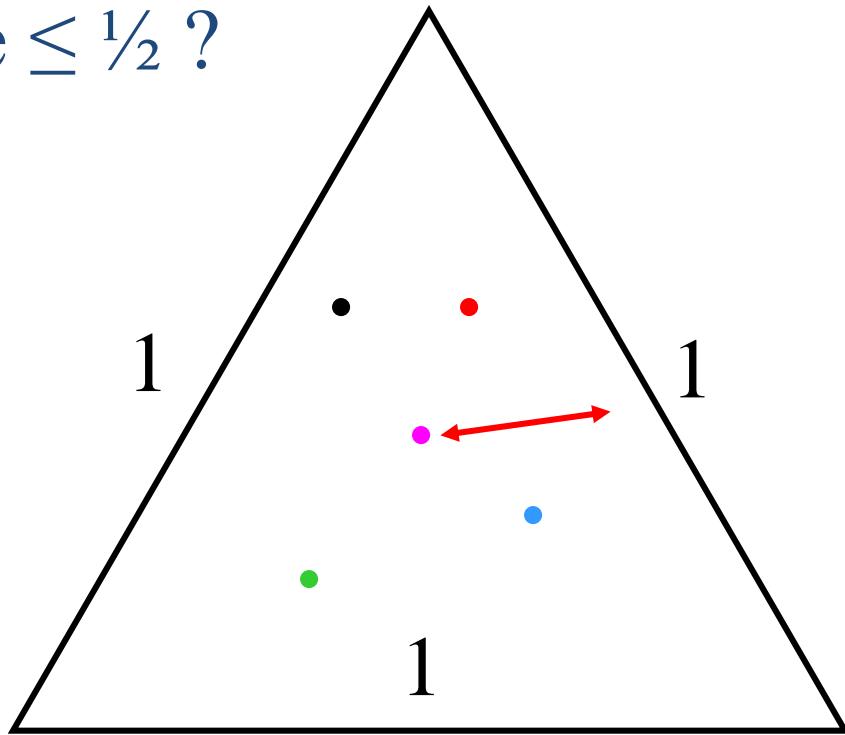
- Useful in analyses and bounds

Problem: Given any five points in/on the unit square, is there always a pair with distance $\leq \frac{1}{\sqrt{2}}$?



- What approaches fail?
- What techniques work and why?
- Lessons and generalizations

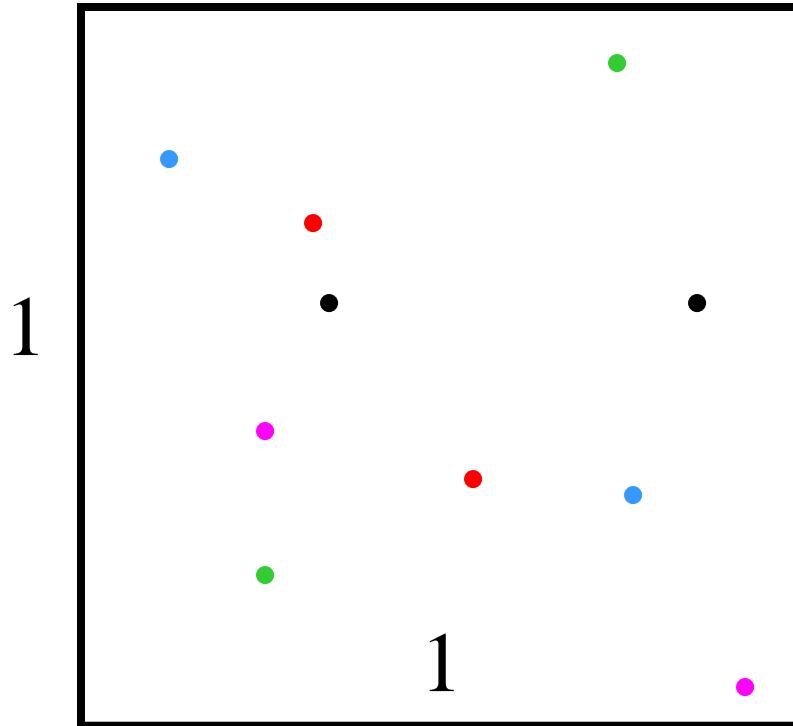
Problem: Given any five points in/on the unit equilateral triangle, is there always a pair with distance $\leq \frac{1}{2}$?



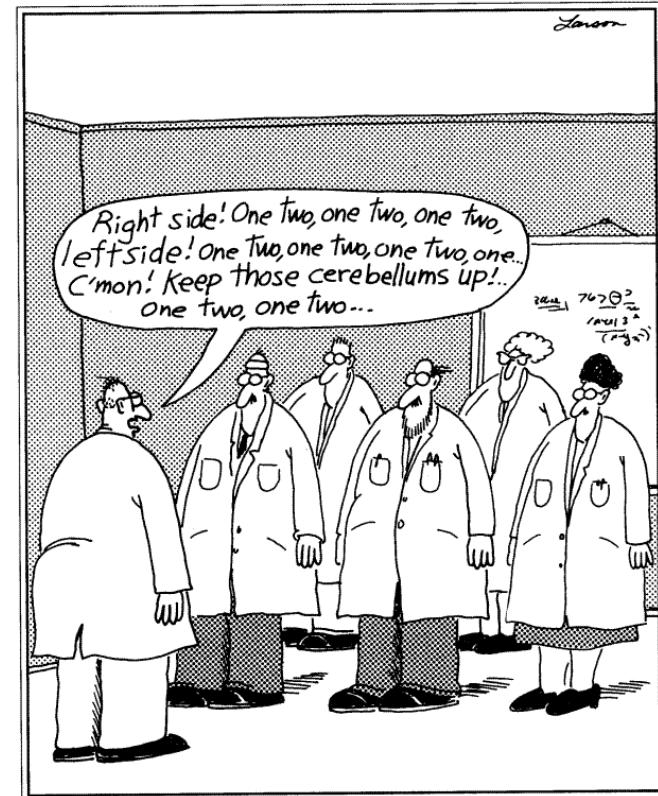
- What approaches fail?
- What techniques work and why?
- Lessons and generalizations



Problem: Given any ten points in/on the unit square, what is the maximum pairwise distance?



- What approaches fail?
- What techniques work and why?
- Lessons and generalizations

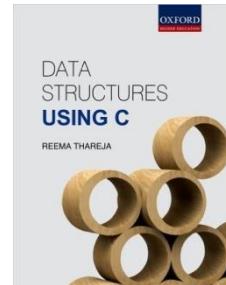
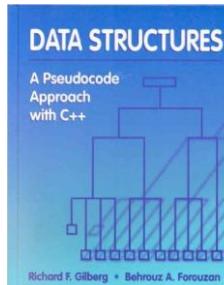
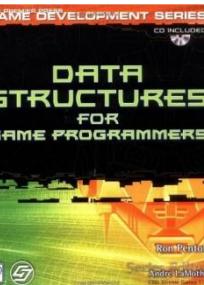
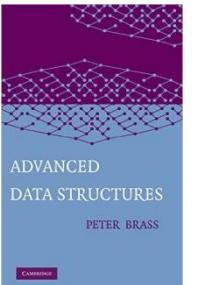
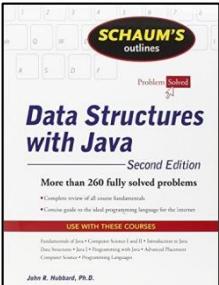
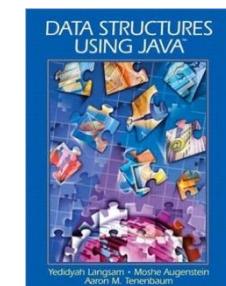
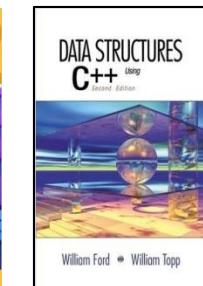
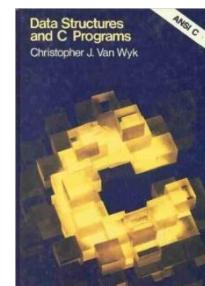
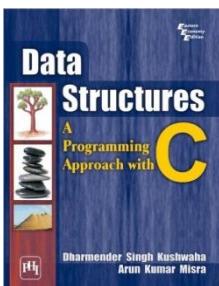
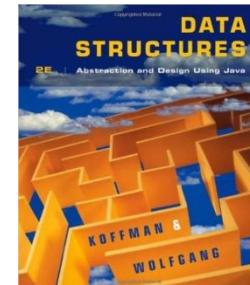
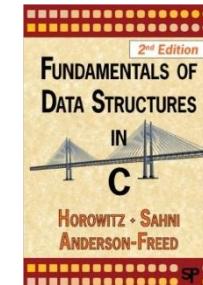
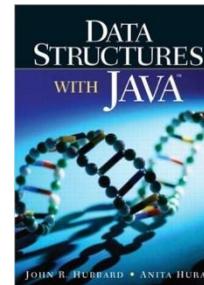
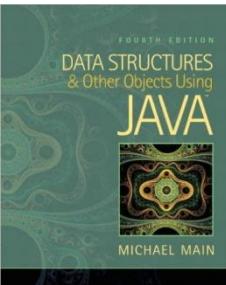
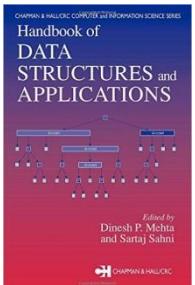
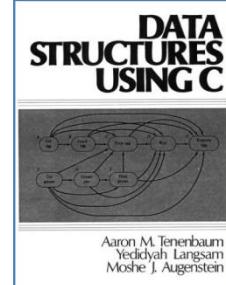
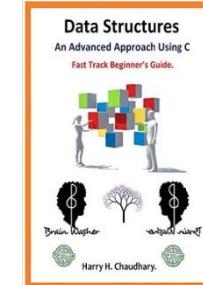
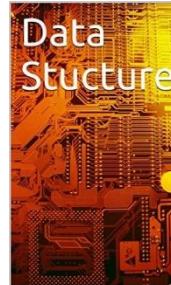
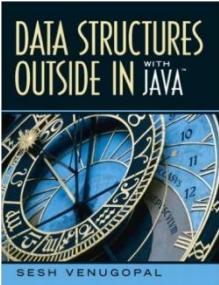


Data Structures



Data Structures

- Techniques for organizing information effectively
- Allowed operations:
 - Initialize
 - Insert
 - Delete
 - Search
 - Min/max
 - Successor
 - Predecessor
 - Merge
 - Split
 - Revert



Data Structures

Primitive types:

- 1. Boolean
- 2. Character
- 3. Floating-point
- 4. Double
- 5. Integer
- 6. Enumerated type

Abstract data types:

- 7. Array
- 8. Container
- 9. Map
- 10. Associative array
- 11. Dictionary
- 12. Multimap
- 13. List
- 14. Set
- 15. Multiset / Bag
- 16. Priority queue
- 17. Queue
- 18. Double-ended queue
- 19. Stack
- 20. String
- 21. Tree
- 22. Graph

Composite types:

- 23. Array
- 24. Record
- 25. Union
- 26. Tagged union

Arrays:

- 27. Bit array
- 28. Bit field
- 29. Bitboard
- 30. Bitmap
- 31. Circular buffer
- 32. Control table
- 33. Image
- 34. Dynamic array
- 35. Gap buffer
- 36. Hashed array tree
- 37. Heightmap
- 38. Lookup table
- 39. Matrix
- 40. Parallel array
- 41. Sorted array
- 42. Sparse array
- 43. Sparse matrix
- 44. Iliffe vector
- 45. Variable-length array

Lists:

- 46. Doubly linked list
- 47. Array list
- 48. Linked list
- 49. Self-organizing list
- 50. Skip list
- 51. Unrolled linked list
- 52. VList
- 53. Xor linked list
- 54. Zipper
- 55. Doubly connected edge list
- 56. Difference list
- 57. Free list

Binary trees:

- 58. AA tree
- 59. AVL tree
- 60. Binary search tree
- 61. Binary tree
- 62. Cartesian tree
- 63. Order statistic tree
- 64. Pagoda
- 65. Randomized binary search tree
- 66. Red-black tree
- 67. Rope

Data Structures

Binary trees (continued):

- 68. Scapegoat tree
- 69. Self-balancing search tree
- 70. Splay tree
- 71. T-tree
- 72. Tango tree
- 73. Threaded binary tree
- 74. Top tree
- 75. Treap
- 76. Weight-balanced tree
- 77. Binary data structure

Trees:

- 78. Trie
- 79. Radix tree
- 80. Suffix tree
- 81. Suffix array
- 82. Compressed suffix array
- 83. FM-index
- 84. Generalised suffix tree
- 85. B-trie
- 86. Judy array
- 87. X-fast trie
- 88. Y-fast trie
- 89. Ctrie

B-trees:

- 90. B-tree
- 91. B+ tree
- 92. B*-tree
- 93. B sharp tree
- 94. Dancing tree
- 95. 2-3 tree
- 96. 2-3-4 tree
- 97. Queap
- 98. Fusion tree
- 99. Bx-tree
- 100. AList

Heaps:

- 101. Heap
- 102. Binary heap
- 103. Weak heap
- 104. Binomial heap
- 105. Fibonacci heap
- 106. AF-heap
- 107. Leonardo Heap
- 108. 2-3 heap
- 109. Soft heap
- 110. Pairing heap
- 111. Leftist heap
- 112. Treap

113. Beap

114. Skew heap

115. Ternary heap

116. D-ary heap

117. Brodal queue

Multiway trees:

- 118. Ternary tree
- 119. K-ary tree
- 120. And-or tree
- 121. (a,b)-tree
- 122. Link/cut tree
- 123. SPQR-tree
- 124. Spaghetti stack
- 125. Disjoint-set data structure
- 126. Fusion tree
- 127. Enfilade
- 128. Exponential tree
- 129. Fenwick tree
- 130. Van Emde Boas tree
- 131. Rose tree

Space-partitioning trees:

- 132. Segment tree
- 133. Interval tree
- 134. Range tree

Data Structures

Space-partitioning trees (cont):

- 135. Bin
- 136. Kd-tree**
- 137. Implicit kd-tree
- 138. Min/max kd-tree
- 139. Adaptive k-d tree
- 140. Quadtree
- 141. Octree
- 142. Linear octree
- 143. Z-order
- 144. UB-tree
- 145. R-tree
- 146. R+ tree
- 147. R* tree
- 148. Hilbert R-tree
- 149. X-tree
- 150. Metric tree
- 151. Cover tree
- 152. M-tree
- 153. VP-tree
- 154. BK-tree
- 155. Bounding interval hierarchy
- 156. BSP tree
- 157. Rapidly exploring random tree

Application-specific trees:

- 158. Abstract syntax tree
- 159. Parse tree**
- 160. Decision tree**
- 161. Alternating decision tree
- 162. Minimax tree**
- 163. Expectiminimax tree
- 164. Finger tree
- 165. Expression tree
- 166. Log-structured merge-tree

Hashes:

- 167. Bloom filter
- 168. Count-Min sketch
- 169. Distributed hash table
- 170. Double Hashing**
- 171. Dynamic perfect hash table
- 172. Hash array mapped trie
- 173. Hash list**
- 174. Hash table**
- 175. Hash tree
- 176. Hash trie
- 177. Koordie
- 178. Prefix hash tree
- 179. Rolling hash

180. MinHash

- 181. Quotient filter
- 182. Ctrie

Graphs:

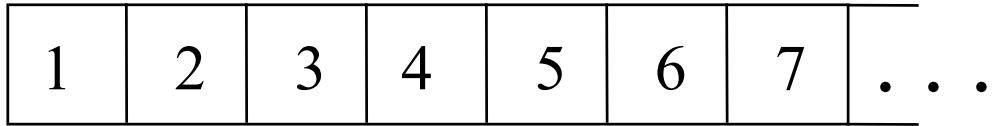
- 183. Graph**
- 184. Adjacency list**
- 185. Adjacency matrix**
- 186. Graph-structured stack
- 187. Scene graph
- 188. Binary decision diagram
- 189. 0-suppressed decision diagram
- 190. And-inverter graph
- 191. Directed graph**
- 192. Directed acyclic graph**
- 193. Propositional dir. acyclic graph
- 194. Multigraph
- 195. Hypergraph

Other:

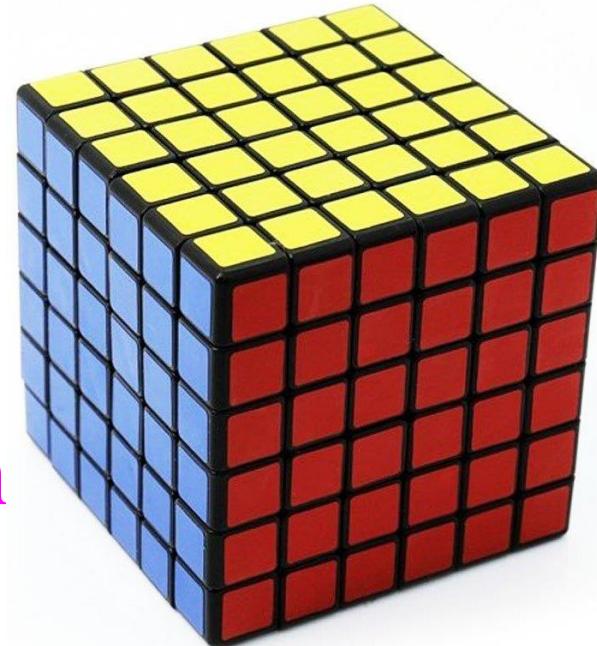
- 196. Lightmap
- 197. Winged edge
- 198. Doubly connected edge list
- 199. Quad-edge
- 200. Routing table
- 201. Symbol table

Arrays

- Sequence of "indexible" locations

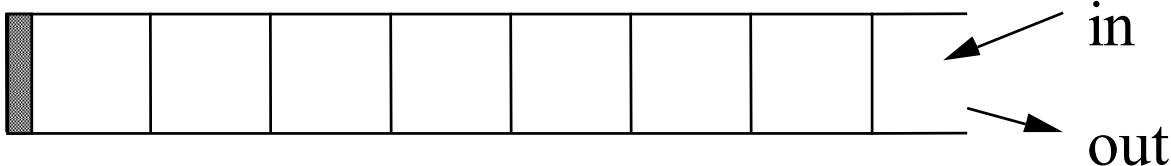


- Unordered:
 - $O(1)$ to add
 - $O(n)$ to search / delete
 - $O(n)$ for min / max
- Ordered:
 - $O(n)$ to add / delete
 - $O(\log n)$ to (binary) search
 - $O(1)$ for min / max

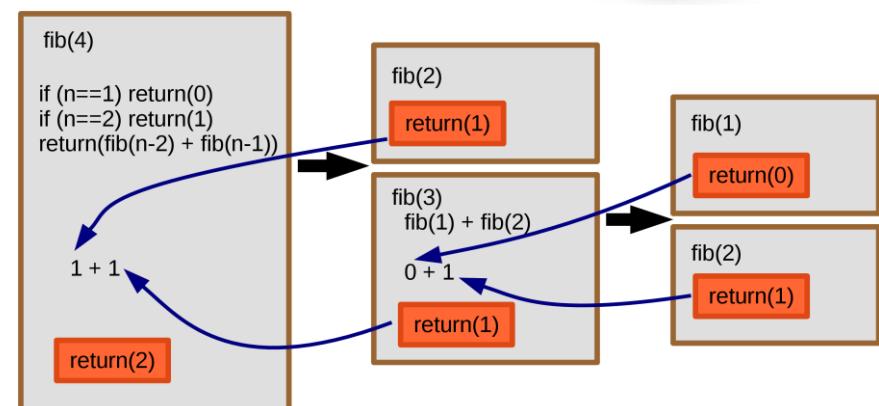


Stacks

- LIFO (Last-In First-Out)

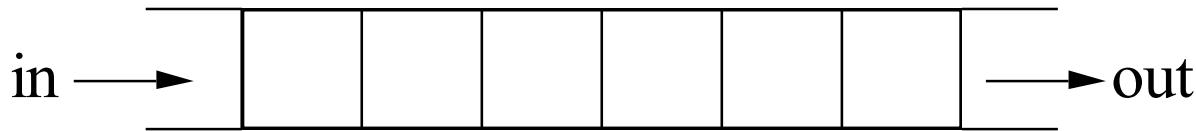


- Operations: push / pop O(1) each
- Can not access “middle”
- Analogy: trays/plates at cafeteria
- Applications:
 - Recursion
 - Compiling / parsing
 - Dynamic binding
 - Web surfing

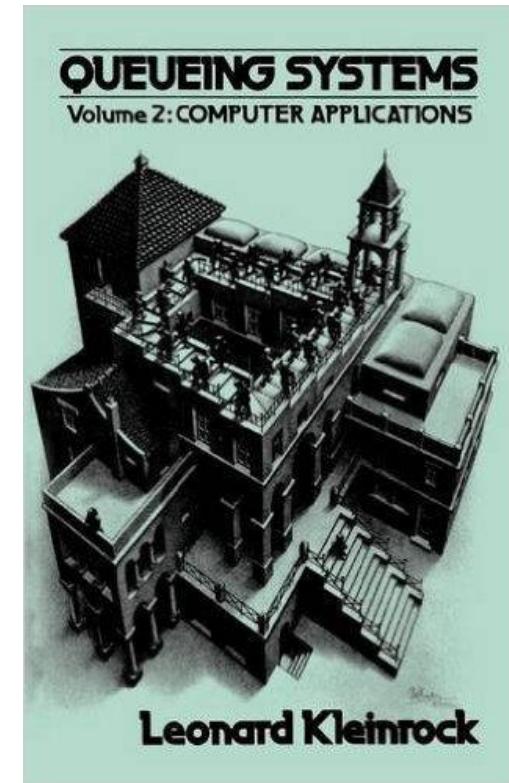


Queues

- **FIFO** (First-In First-Out)



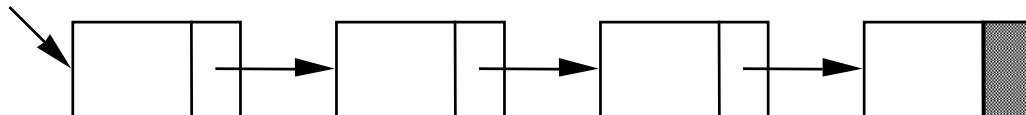
- Operations: **push** / **pop** O(1) each
- Can not access “middle”
- Analogy: line at the store
- Applications:
 - Simulations
 - Scheduling
 - Networks
 - Operating systems



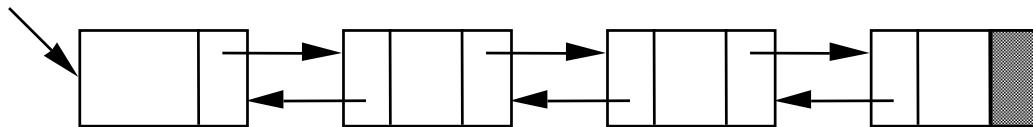
Linked Lists

- Successor / predecessor pointers
- Types:

- Single linked



- Double linked



- Circular

- Operations:

- Add: $O(1)$ time

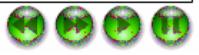
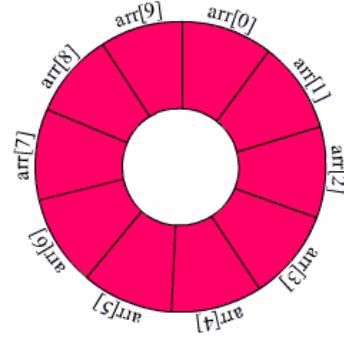
- Search: $O(n)$ time

- Delete: $O(1)$ time (given pointer)

Building a linked list

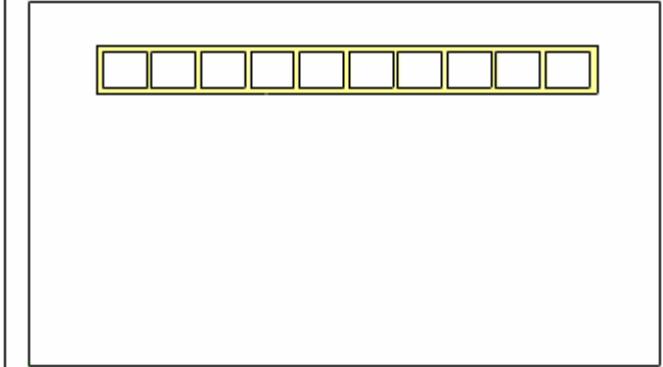
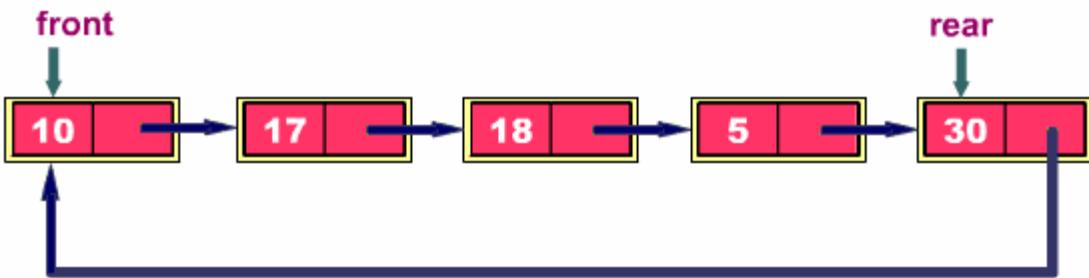


Building of Circular Queue

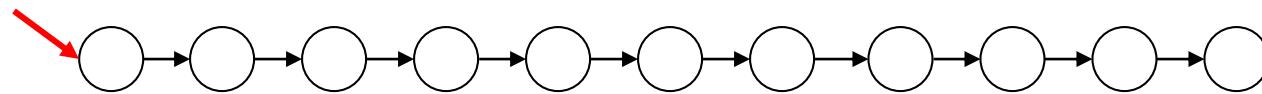


Circular queue as a linked list

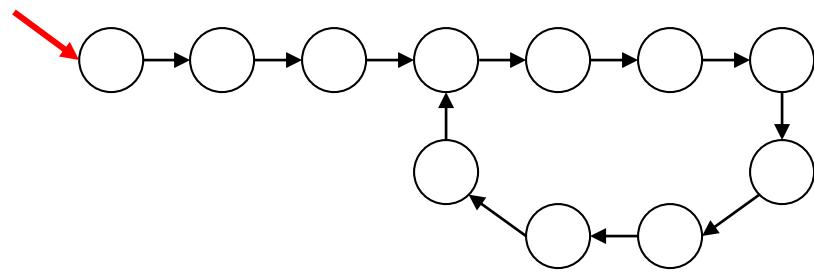
Deletion of a node



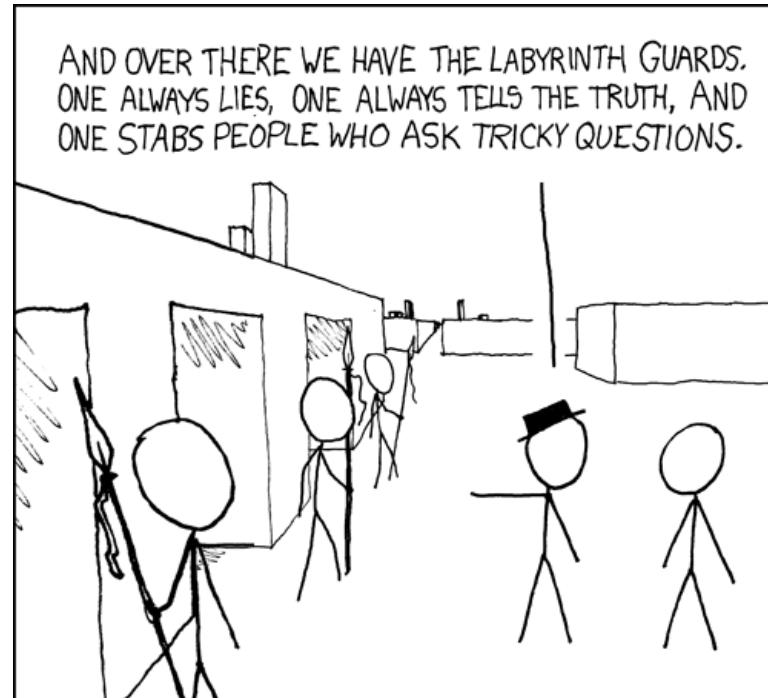
Extra Credit Problem: Given a pointer to a **read-only** (unmodifiable) linked list containing an **unknown number** of nodes n , devise an **$O(n)$ -time** and **$O(1)$** space algorithm that determines whether that list contains a cycle.



or

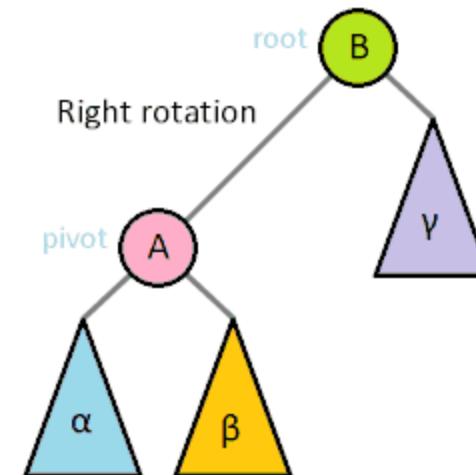
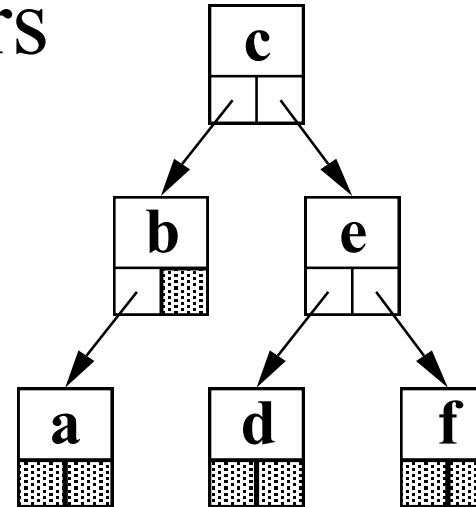


- What approaches fail?
- What techniques work and why?
- Lessons and generalizations



Trees

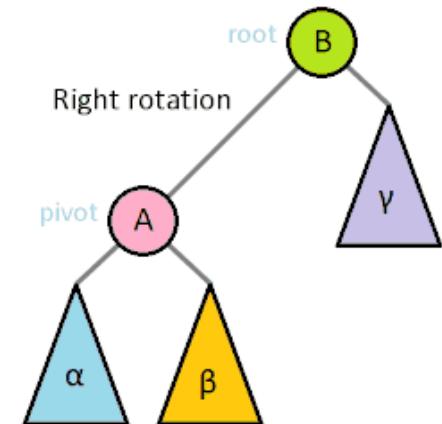
- Parent/children pointers
- Binary / N-ary
- Ordered / unordered
- Height-balanced:
 - B-trees
 - AVL trees
 - Red-black trees
 - 2-3 trees
 - **add / delete / search** in $O(\log n)$ time



B-Trees



- Multi-rotations occur infrequently
- Rotations don't propagate far
- Larger tree \Rightarrow fewer rotations
- Same for other height-balanced trees
- Non-balanced search trees **average** $O(\log n)$ height

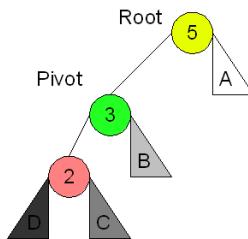


Height-Balanced AVL Trees

There are 4 cases in all, choosing which one is made by seeing the direction of the first 2 nodes from the unbalanced node to the newly inserted node and matching them to the top most row.

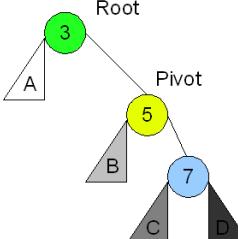
Root is the initial parent before a rotation and **Pivot** is the child to take the root's place.

Left Left Case



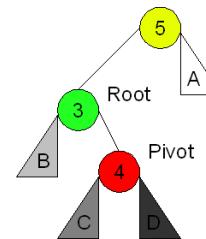
Right Rotation

Right Right Case



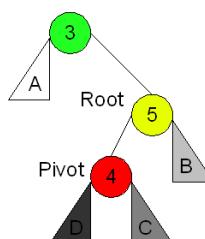
Left Rotation

Left Right Case

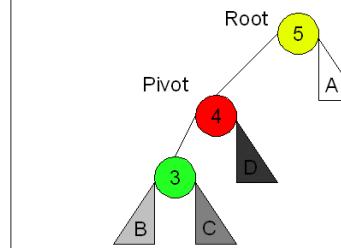
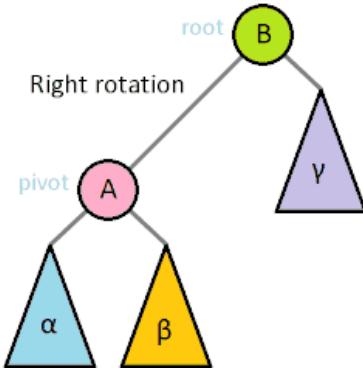


Left Rotation

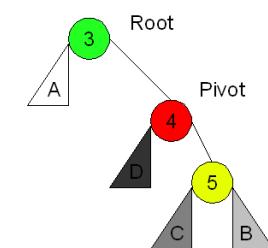
Right Left Case



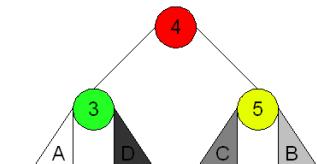
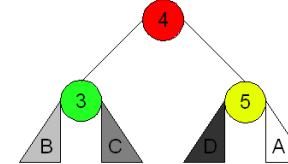
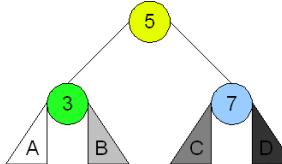
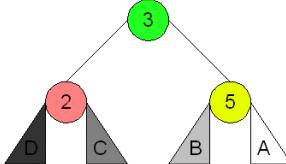
Right Rotation



Right Rotation

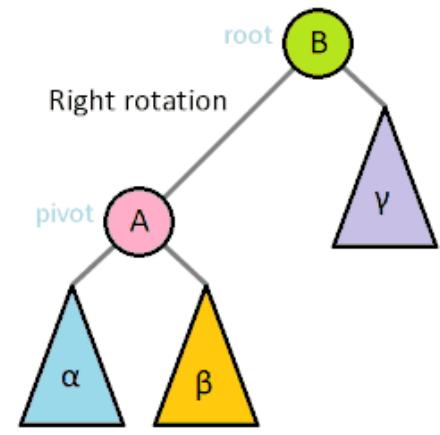


Left Rotation

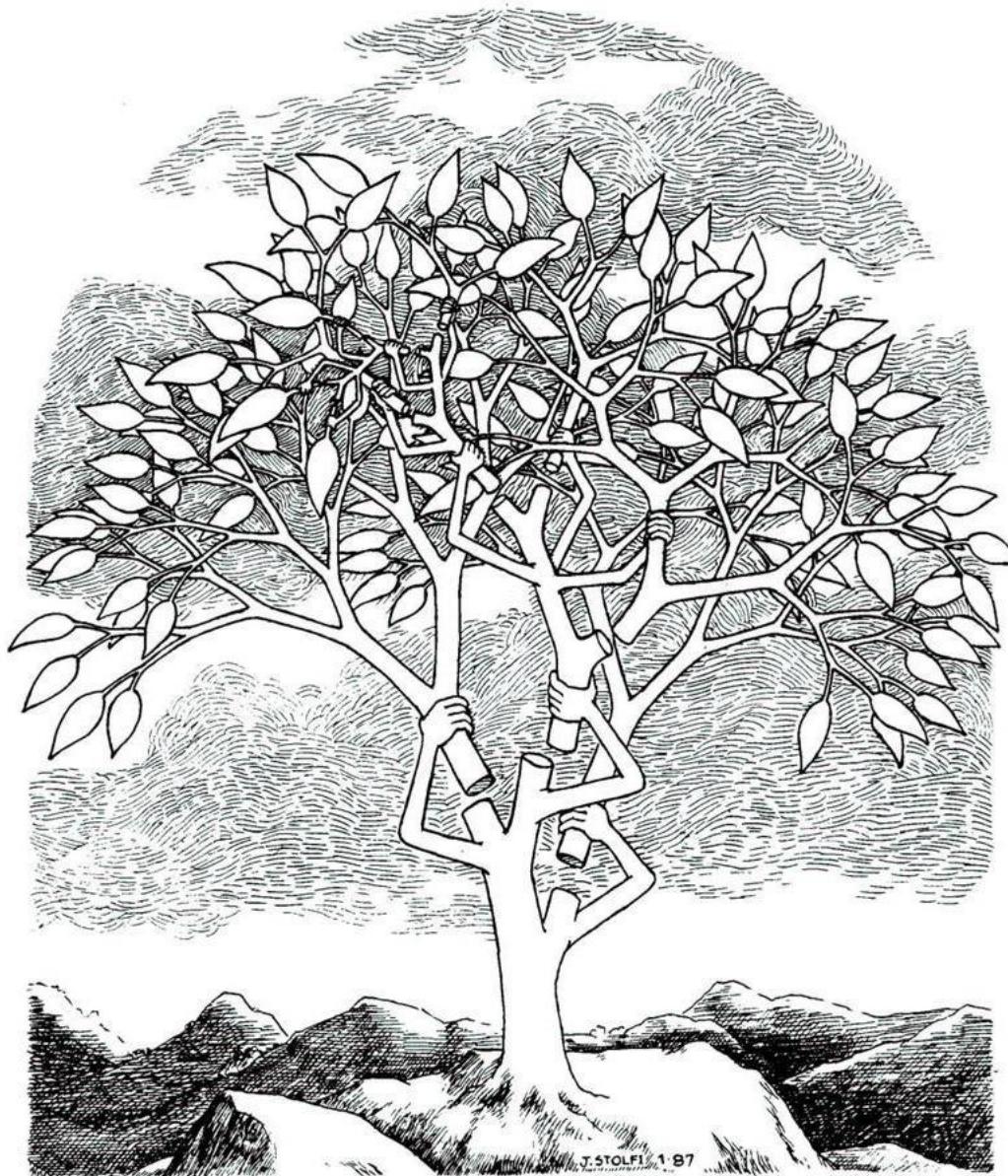
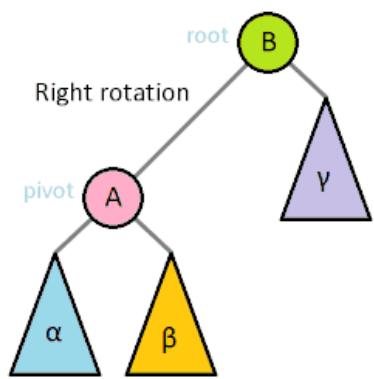


AVL Trees

- Multi-rotations occur infrequently
- Rotations don't propagate far
- Larger tree \Rightarrow fewer rotations
- Same for other height-balanced trees
- Non-balanced trees **average $O(\log n)$** height



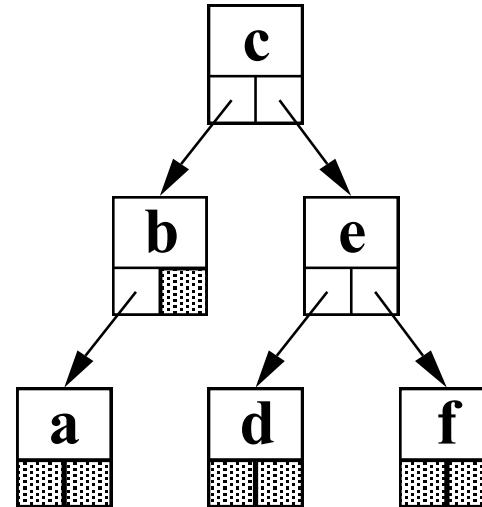
Self-Adjusting Trees



Tree Traversal

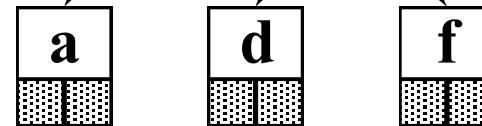
- Pre-order:

1. process node
 2. visit children
- ⇒ c b a e d f



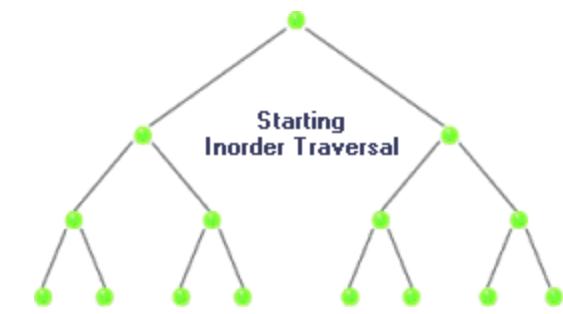
- Post-order:

1. visit children
 2. process node
- ⇒ a b d f e c



- In-order:

1. visit left-child
 2. process node
 3. visit right-child
- ⇒ a b c d e f



Heaps

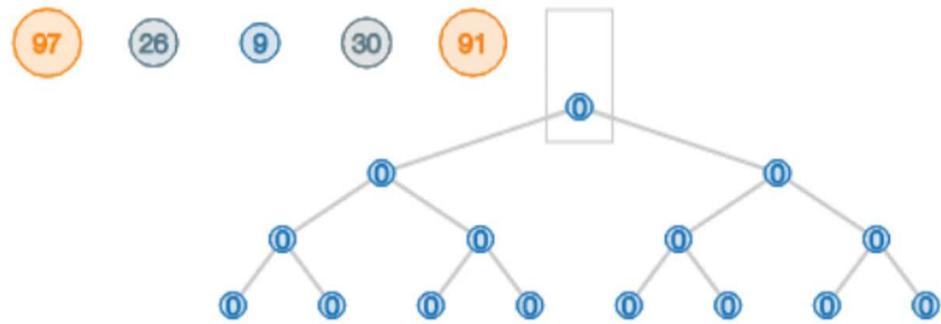
- A tree where all of each node's children have larger / smaller “keys”

- Can be implemented using binary tree or array

- Operations:

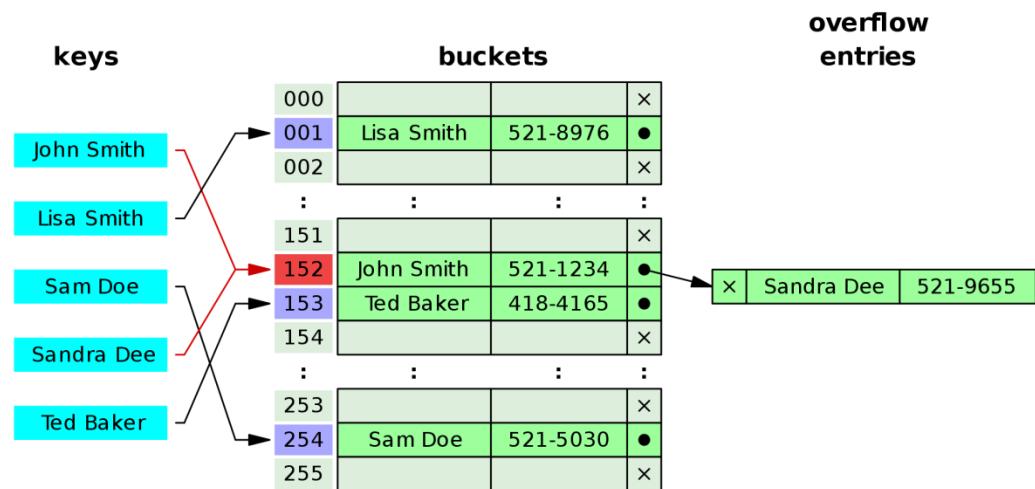
- Find min/max: $O(1)$ time
- Add: $O(\log n)$ time
- Delete: $O(\log n)$ time
- Search: $O(n)$ time

6 5 3 1 8 7 2 4



Hash Tables

- Direct access
- Hash function
- Collision resolution:
 - Chaining
 - Linear probing
 - Double hashing
- Universal hashing
- $O(1)$ average access
- $O(n)$ worst-case access



Q: Improve **worst-case** access to $O(\log n)$?

3.4 LINEAR PROBING HASH TABLE DEMO



click to begin demo